4-CYCLE PROPERTIES FOR CHARACTERIZING RECTAGRAPHS AND HYPERCUBES

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Abstract. A (0, 2)-graph is a connected graph, where each pair of vertices has either 0 or 2 common neighbours. These graphs constitute a subclass of $(0, \lambda)$ -graphs introduced by Mulder in 1979. A rectagraph, well known in diagram geometry, is a triangle-free (0, 2)-graph. (0, 2)-graphs include hypercubes, folded cube graphs and some particular graphs such as icosahedral graph, Shrikhande graph, Klein graph, Gewirtz graph, etc. In this paper, we give some local properties of 4-cycles in $(0, \lambda)$ -graphs and more specifically in (0, 2)-graphs, leading to new characterizations of rectagraphs and hypercubes.

Keywords: hypercube; (0, 2)-graph; rectagraph; 4-cycle; characterization

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1. INTRODUCTION

A (0, 2)-graph is a connected graph, where each pair of vertices has either 0 or 2 common neighbours. This class of graphs is a special case of $(0, \lambda)$ -graphs, introduced by Mulder, who showed that they are regular in [6]. A rectagraph, well-known in diagram geometry, see [9], is a (0, 2)-graph without any triangles.

Since every pair of adjacent edges engenders a unique 4-cycle, we can easily compute the number of 4-cycles in a (0, 2)-graph. In fact, there are |V(G)|d(d-1)/8 4-cycles in a (0, 2)-graph G of order |V(G)| and degree d.

(0, 2)-graphs were studied in various contexts; existence and construction of (0, 2)-graphs were intensively studied by several researchers, we can cite for instance [2], [3]. Another important aspect is characterizing hypercubes as (0, 2)-graphs; due to their remarkable properties and multiple applications, many authors have investigated this topic in order to give a new point of view which can be used for recognizing and constructing hypercubes, see [1], [5], [6], [7], [10].

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In this paper, we give new characterizations of hypercubes and rectagraphs, in addition to many local properties of 4-cycles in $(0, \lambda)$ -graphs and more specifically in (0, 2)-graphs, using an edge binary relation introduced by Sabidussi in his work [11] on Cartesian graph products.

2. Definitions and notation

Let G be a simple connected graph with V(G) its vertex set and E(G) its edge set, let x and y be two vertices in V(G). We denote by N(x) the set of neighbours of x and by N(x, y) the set of common neighbours of x and y. An edge $\{x, y\}$ will be denoted xy, where vertices x and y are its *extremities*.

A hypercube or a d-dimensional cube is the graph, denoted Q_d , with $V(Q_d) = \{0,1\}^d$ where the d-tuples x and y are joined if and only if they differ in exactly one coordinate. Let G be a graph and let $u, v \in V(G)$. The distance d(u, v) between two vertices u and v is the length of a shortest (u, v)-path. The diameter of the graph G, denoted diam(G), is the maximum distance between any pairs of vertices. The interval I(u, v) can be defined by $I(u, v) = \{w \in V(G): d(u, v) = d(u, w) + d(w, v)\}$.

A vertex \bar{u} is an *antipode* of a vertex u if $I(u, \bar{u}) = V(G)$. A graph G is an *antipodal graph* if each vertex has a unique antipode. A subset $W \subseteq V(G)$ is *convex* if $I(u, v) \subseteq W$ for any $u, v \in W$. A graph G is *interval monotone* if each of its intervals is convex. A graph G is *interval regular* if $|I(u, v) \cap N(u)| = d(u, v)$ for any two vertices u and v in V(G). Note that a hypercube is characterized as a (0, 2)-graph, see [5], [6], interval monotone [4], interval regular [8] and an antipodal graph [10].

Let G be a graph. A 4-cycle denoted C is defined by the sequence of its four vertices. The set of all 4-cycles in G is denoted $\Phi(G)$.

Let $e, e \in E(G)$, we define a binary relation θ on E(G), where $e\theta e$ if there is a 4-cycle $C \in \Phi(G)$ such that e, e are nonadjacent edges of C. We say that the edge e is *parallel* to e and denote $\theta(e) = \{e \in E(G): e\theta e\}$. Note that θ was introduced by Sabidussi in [11] as the relation $\sim^{(0)}$.

Let G be a (0,2)-graph and let e be an edge in E(G). We define the edge level decomposition relative to e by defining the subsets $N_i(e) \subseteq E(G)$, $0 \leq i \leq m$, which constitute the levels of the decomposition, where $N_0(e) = \{e\}$ and $e_i \in N_i(e)$, if there is at least one edge $e_{i-1} \in N_{i-1}(e)$ such that $e_i \in \theta(e_{i-1})$ for $1 \leq i \leq m$, where m denotes the index of the last level in the decomposition. It is then clear that in any edge level decomposition relative to e, every edge f = xy in E(G) is either in $N_k(e), 0 \leq k \leq m$, or belongs to $E(G) \setminus \bigcup_{i=1}^m N_i(e)$. In this case, its two extremities x and y are incident with two different edges from $\bigcup_{i=1}^{m} N_i(e)$, which are either in the same level or different levels.

3. Some 4-cycle properties in (0, 2)-graphs

In this section, we give first some local properties of 4-cycles in $(0, \lambda)$ -graphs. After that, other results, using the relation θ , are presented for the specific case $\lambda = 2$, including a characterization of rectagraphs.

Lemma 3.1. If G is a $(0, \lambda)$ -graph, then every vertex in V(G) is in $d(d-1) \times (\lambda - 1)/2$ 4-cycles.

Proof. Let u be a vertex. Since G is regular, |N(u)| = d and for each pair of vertices $v, w \in N(u)$, there are $(\lambda - 1)$ 4-cycles containing u, v, w at one time. Thus the number of 4-cycles containing u equals $(\lambda - 1)\binom{d}{2}$.

Lemma 3.2. If G is a $(0, \lambda)$ -graph, then every edge e is contained in $(\lambda - 1)(d - 1)$ 4-cycles.

Proof. Let e = uv be an edge. For each edge e_u incident to u, there are $\lambda - 1$ edges $e_1, e_2, \ldots, e_{\lambda-1}$ incident to v such that e, e_u and $e_i, 1 \leq i \leq \lambda - 1$ are contained in a 4-cycle. Since G is d-regular, e is then contained in $(\lambda - 1)(d - 1)$ 4-cycles. \Box

Lemma 3.3. Let G be a $(0, \lambda)$ -graph. If there is an edge $e \in E(G)$ such that $\theta(e)$ is a nonempty matching, then $\lambda = 2$.

Proof. Let G be a $(0, \lambda)$ -graph. It is easy to see that if $\lambda = 1$ then $\Phi(G) = \emptyset$. Therefore, $\theta(e) = \emptyset$ for all edges in E(G). Note that if $\lambda \ge 2$ and $|E(G)| \ge 2$, then every two adjacent edges are contained in $(\lambda - 1)$ 4-cycles, thus for any edge e, there are at least $\lambda - 1$ adjacent edges that are parallel to e. Therefore, if there is an edge e such that $\theta(e)$ is a matching, then we have necessarily $\lambda = 2$.

Lemma 3.4. In a (0, 2)-graph G, we have

$$d-2 \leq |\theta(e)| \leq d-1, \quad e \in E(G).$$

Furthermore, G is K_4 free if and only if $|\theta(e)| = d - 1$ for all $e \in E(G)$.

Proof. Let G be a (0,2)-graph. First, we prove that $d-2 \leq |\theta(e)| \leq d-1$ for all $e \in E(G)$.

Let e = uv be an edge in E(G). According to Lemma 3.2, e is contained in (d-1)4-cycles, and by definition of the relation θ , we can easily deduce that $|\theta(e)| \leq d-1$. On the other hand, assume that $|\theta(e)| \leq d-3$, then there are $e_1 = xy$, $e_2 = zt \in \theta(e)$, $e_1 \neq e_2$, such that e and e_1 are contained in two 4-cycles C_1 , C_2 and e and e_2 are contained in two other 4-cycles C_3 , C_4 . Thus on one hand we have $N(u, v) = \{x, y\}$, and on the other hand we have $N(u, v) = \{z, t\}$. By the (0, 2)-property, this implies that $\{x, y\} = \{z, t\}$ which contradicts the fact that $e_1 \neq e_2$.

Note that if the edge $e = uv \in E(G)$ is such that $|\theta(e)| < d - 1$, then there is an edge $\bar{e} \in \theta(e)$ such that e and \bar{e} are contained in two different 4-cycles C_1, C_2 , with $C_1 = uvwt$ and $C_2 = uvtw$, which means that vertices u, v, w, and t are pairwise adjacent and the subgraph induced by the set $\{u, v, w, t\}$ is isomorphic to K_4 .

Conversely, if the subgraph induced by $Y = \{u, v, w, t\} \subseteq V(G)$ is isomorphic to K_4 , then there are two 4-cycles $C_1, C_2 \in \Phi(G)$ containing e and \bar{e} with $C_1 = uvwt$ and $C_2 = uvtw$, hence $|\theta(e)| = d - 2 < d - 1$.

Lemma 3.5. Let G be a (0,2)-graph and let $e \in E(G)$. If $\bar{e} \in \theta(e)$, then there is at most one edge $\bar{e} \in \theta(e)$ such that \bar{e}, \bar{e} are adjacent. In this case, G contains $K_4 - e$ as an induced subgraph.

Proof. Let G be a (0,2)-graph. Let $\dot{e} = uv \in E(G)$ and $\bar{e} = wz \in \theta(\dot{e})$. Suppose that there are two edges e_1, e_2 adjacent to \bar{e} such that $e_1, e_2 \in \theta(\dot{e})$. We have then, without loss of generality, two cases:

Case 1: $e_1 = wt$ and $e_2 = wy$. Since $\bar{e} \in \theta(\acute{e})$, there is a 4-cycle $C_1 = uvwz$ and $N(u,w) = \{v,z\}$. On the other hand, $e_1 = wt \in \theta(\acute{e})$ implies that $ut \notin E(G)$ (otherwise $N(u,w) = \{v,z,t\}$) thus there is a 4-cycle $C_2 = uvtw$ and $N(v,w) = \{u,t\}$. Since $e_2 = wy \in \theta(\acute{e})$, there is a 4-cycle C_3 such that $C_3 = uvyw$ or $C_3 = uvwy$. But if $C_3 = uvyw$, then $N(v,w) = \{u,t,y\}$ and if $C_3 = uvwy$, then $N(u,w) = \{v,z,y\}$, which is impossible in both the cases.

Case 2: $e_1 = wt$ and $e_2 = zy$. Since $\bar{e} \in \theta(\acute{e})$, there is a 4-cycle $C_1 = uvwz$ and $N(v, z) = \{u, w\}$. We have also $e_1 = wt \in \theta(\acute{e})$ and $ut \notin E(G)$ (otherwise $N(u, w) = \{v, z, t\}$) which implies that there is $C_2 = uvtw$ and thus $N(v, w) = \{u, t\}$. On the other hand, $e_2 = zy \in \theta(\acute{e})$ implies that there is $C_3 \in \Phi(G)$ such that $C_3 = uvyz$ or $C_3 = uvzy$. But if $C_3 = uvyz$, then $N(v, z) = \{u, w, y\}$, and if $C_3 = uvzy$, then $N(v, w) = \{u, z, t\}$, which is also impossible.

Note that if $\bar{e} = wz$ and $e_1 = wt$ are parallel to $\dot{e} = uv$, then there are two 4-cycles $C_1 = uvwz$ and $C_2 = uvtw$ in $\Phi(G)$ and the subgraph induced by $\{u, v, z, w\}$ is isomorphic to $K_4 - e$ with e = uz, which concludes this proof.

Theorem 3.6. Let G be a $(0, \lambda)$ -graph of degree d. G is a rectagraph if and only if $\theta(\acute{e})$ is a matching of (d-1) edges for any edge \acute{e} .

Proof. Let G be a rectagraph and let $e \in E(G)$. Since G is triangle free, G is also K_4 free and $(K_4 - e)$ free. According to Lemma 3.4, $|\theta(e)| = d - 1$ and by Lemma 3.5, there are no adjacent edges in $\theta(e)$. It follows that $\theta(e)$ is a matching of d-1 edges.

Conversely, if $\theta(\acute{e})$ is a matching of d-1 edges for every edge $\acute{e} \in E(G)$ then according to Lemma 3.3, G is a (0,2)-graph. Now assume that G is a (0,2)-graph that contains triangles, and let x, y, z be three vertices inducing a triangle. Let $e_1 = xy, e_2 = yz$ and $e_3 = xz$. Since $z \in N(x, y)$, there is a unique vertex t such that $N(x, y) = \{t, z\}$. Hence there is $C_1 \in \Phi(G)$ such that $C_1 = xzyt$ and it follows that $\hat{e} = xt \in \theta(e_2)$. On the other hand, since $y \in N(x, z)$, there is a unique vertex w such that $N(x, z) = \{w, y\}$. Note that $w \neq t$, otherwise the subgraph induced by $\{x, y, z, t\}$ is isomorphic to K_4 and according to Lemma 3.4, $|\theta(\hat{e})| = d - 2 < d - 1$. Consequently, there is a 4-cycle $C_2 \in \Phi(G)$ such that $C_2 = xyzw$ and $\tilde{e} = xw \in \theta(e_2)$. This means that $\theta(e_2)$ is not a matching, since \hat{e} and \tilde{e} are adjacent, which contradicts our assumption.

4. New characterizations of hypercubes

Hypercubes constitute a remarkable class of graphs with very interesting properties, including the (0, 2)-property. In this section, we show new characterizations of a hypercube as a (0, 2)-graph. But first, let us recall important results due to Mulder and Laborde and Rao Hebbare that we shall use:

Proposition 4.1 (Mulder [6], Laborde and Rao Hebbare [5]). Let G be a (0, 2)-graph of degree d, then $|V(G)| \leq 2^d$. Furthermore, G is a hypercube of dimension d if and only if $|V(G)| = 2^d$.

Mulder has also shown the following result:

Proposition 4.2 (Mulder [8]). Let G be a connected graph. G is a hypercube if and only if G is bipartite and interval regular.

Theorem 4.3. Every interval monotone rectagraph is interval regular.

Proof. Let G be an interval monotone rectagraph. The proof is by induction on d(u, v) where $u, v \in V(G)$. If d(u, v) = 2, then $I(u, v) = \{x, y\}$ and $|N(u) \cap I(u, v)| = |N(u, v)| = 2$. Let $u, v \in V(G)$ such that $d(u, v) = k \ge 3$, and let $w \in N(u) \cap I(u, v)$.

By induction hypothesis, we have $|N(w) \cap I(w, v)| = d(w, v) = k - 1$. We denote by $t_1, t_2, \ldots, t_{k-1}$, the neighbours of w in I(w, v). Since $d(u, t_i) = 2$, $1 \leq i \leq k - 1$, there are $w_1, w_2, \ldots, w_{k-1} \in V(G)$ with $w_i \neq w$, $1 \leq i \leq k - 1$, such that $N(u, t_i) = \{w, w_i\}$. Hence $d(w_i, v) = d(t_i, v) + 1 = k - 1$ and $w_i \in N(u) \cap I(u, v)$. Note that $w_i \neq w_j$ for $i \neq j$, otherwise $N(w, w_i) = N(w, w_j) = \{u, t_i, t_j\}$. Assume now that there is a vertex $x \in N(u) \cap I(u, v)$ such that $x \neq w, w_1, w_2, \ldots, w_{k-1}$. Therefore $u \in N(x, w)$ and there is a unique vertex $y \neq u$ such that $y \in N(x, w)$. Note that $y \neq t_i, 1 \leq i \leq k - 1$, otherwise $N(u, t_i) = \{x, w, w_i\}$. Since G is interval monotone and triangle free, $y \in I(x, w) \subset I(u, v)$, and consequently, $y \in N(w) \cap I(u, v)$, which contradicts the induction hypothesis.

Corollary 4.4. Let G be a graph. G is a hypercube if and only if G is a bipartite interval monotone rectagraph.

Proof. A hypercube is a bipartite interval monotone rectagraph. Conversely, according to Theorem 4.3, an interval monotone rectagraph G is interval regular. Since G is bipartite, G is a hypercube by Proposition 4.2.

Theorem 4.5. Let G be a rectagraph of degree d. G is a hypercube of dimension d if and only if in an edge level decomposition relative to any edge e, we have $|N_0(e)| =$ $|N_m(e)| = 1$, $|N_{i-1}(e) \cap \theta(e_i)| = i$ and $|N_{i+1}(e) \cap \theta(e_i)| = d - 1 - i$ for any edge $e_i \in N_i(e)$, $1 \le i \le m$.

Proof. Let Q_d be a hypercube of dimension d, let e = uw be an edge in $E(Q_d)$ and let $N_0(e), N_1(e), \ldots, N_m(e)$ be the subsets of edges in the level decomposition relative to the edge e.

First, we prove that $|N_m(e)| = 1$ and m = d - 1. Let v be the unique antipode of u in Q_d and t the unique antipode of w, thus $I(u, v) = I(w, t) = V(Q_d)$ and d(u, v) = d(w, t) = d. Since $v \in I(w, t)$, we have d(w, t) = d = d(w, v) + d(v, t). On the other hand, $w \in I(u, v)$ implies that d(u, v) = d = d(u, w) + d(w, v). We can easily deduce that d(u, w) = d(v, t), and v, t are adjacent. Since d(u, t) = d(w, v) = d - 1, the edge $vt \in N_m(e)$ and m = d - 1. This edge is unique in $N_{d-1}(e)$, otherwise there would be another pair of vertices x and y which are antipodes of u and v, respectively. Hence $|N_m(e)| = 1$.

Now let $e_i = xy \in N_i(e)$. Since d(u, x) = i and Q_d is interval regular, $|N(x) \cap I(u, x)| = i$, thus there are *i* neighbours x_1, x_2, \ldots, x_i of *x* in I(u, x). Note that $d(u, x_j) = i - 1$ and $x \in N(y, x_j), 1 \leq j \leq i$, which implies that there are *i* vertices, y_1, y_2, \ldots, y_i such that $y_j \in N(y, x_j), 1 \leq j \leq i$. It follows that $x_j y_j \in \theta(e_i), 1 \leq j \leq i$ and we have necessarily the edges $x_j y_j \in N_{i-1}(e), 1 \leq j \leq i$. Thus $|N_{i-1}(e) \cap \theta(e_i)| = i$. In the same way, let $vt \in N_{d-1}(e)$ be such that d(u, t) =

d(w, v) = d - 1. So d(x, t) = d - 1 - i and there are d - 1 - i neighbours of the vertex x (say $\dot{x}_1, \dot{x}_2, \dots, \dot{x}_{d-1-i}$) in I(x, t). Since $x \in N(y, \dot{x}_i), 1 \leq j \leq d-1-i$, there are vertices $\dot{y}_1, \dot{y}_2, \ldots, \dot{y}_{d-1-i}$, such that $\dot{y}_j \in N(y, \dot{x}_j), 1 \leq j \leq d-1-i$. Thus $\dot{x}_{i}\dot{y}_{i} \in N_{i+1}(e), \ 1 \leq j \leq d-1-i \text{ and } |N_{i+1}(e) \cap \theta(e_{i})| = d-1-i.$

Conversely, let G be a rectagraph satisfying the hypothesis of Theorem 4.5 and let $e \in E(G)$. Let $N_0(e), N_1(e), \ldots, N_m(e)$ be the subsets of edges in the level decomposition relative to the edge e. Note that for any edge $e_i \in N_i(e)$, every edge of $\theta(e_i)$ is either in $N_{i-1}(e)$ or $N_{i+1}(e)$.

First we prove that $N_i(e)$ is a matching for $1 \leq i \leq m$. So let us assume the contrary, and choose the smallest k such that $N_k(e)$ is not a matching. So there are at least two adjacent edges in $N_k(e)$, say $e_k = xy$ and $\dot{e}_k = yz$. Since $|N_{k-1}(e) \cap \theta(e_k)| =$ k, there is an edge $e_{k-1} = \dot{x}\dot{y}$ parallel to e_k in $N_{k-1}(e)$. Note that $\dot{e}_k \notin \theta(e_{k-1})$, since G is a rectagraph. Therefore, since $y \in N(y, z)$, there is necessarily another vertex $\dot{z} \in N(\dot{y}, z)$, thus the edge $\dot{e}_{k-1} = \dot{y}\dot{z} \in \theta(\dot{e}_k)$ and consequently, $\dot{e}_{k-1} \in N_{k-1}(e)$. But edges e_{k-1} and e_{k-1} are adjacent, which contradicts the fact that k is the smallest integer such that $N_k(e)$ contains adjacent edges. It follows that $N_i(e)$ is a matching for $1 \leq i \leq m$.

Now let us prove, by induction on *i*, that $|N_i(e)| = \binom{d-1}{i}, 0 \leq i \leq m$. For i = 0, we have $|N_0(e)| = 1 = \binom{d-1}{0}$. Assume now that for all $k \leq i - 1$, $|N_k(e)| = \binom{d-1}{k}$. Since each edge $e_{i-1} \in N_{i-1}(e)$ is parallel to d-i edges in $N_i(e)$, the total number of 4-cycles lying between $N_{i-1}(e)$ and $N_i(e)$ equals $\binom{d-1}{i-1}(d-i)$.

On the other hand, each edge $e_i \in N_i(e)$ is parallel to *i* edges from $N_{i-1}(e)$. Thus, the number of 4-cycles lying between $N_i(e)$ and $N_{i-1}(e)$ equals $|N_i(e)|i$, and it follows that

$$|N_i(e)| = \frac{\binom{d-1}{i-1}(d-i)}{i} = \binom{d-1}{i}.$$

Note that $|N_m(e)| = 1 = \binom{d-1}{m}$ implies m = d - 1.

By counting the total number of edges in the levels, we have

$$|N_0(e)| + |N_1(e)| + \ldots + |N_{d-1}(e)| = \sum_{i=0}^{d-1} {d-1 \choose i} = 2^{d-1}.$$

Since each level is a matching, it follows that $|V(G)| = 2 \cdot 2^{d-1} = 2^d$. According to Proposition 4.1, G is a hypercube of dimension d.

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